

Implementation of the Kazhdan–Lusztig algorithm

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Some notes to myself about the implementation of the K-L algorithm. The main reference are David's notes [2] from a Park City conference.

We assume that we are given tables of the cross-actions and Cayley transforms. For $s \in S$, we denote C^s the Cayley transform going from more compact to less compact Cartans, C_s the inverse Cayley transform. We also use γ^s (resp. γ_s^\pm) for $C^s(\gamma)$ when C^s is one- (resp. two-) valued on γ ; similarly γ_s (resp. γ_s^\pm) for C_s .

1 Hecke algebra action

We denote \mathcal{D} the set of parameters (of course the existence of the Cayley transform and cross action tables imply that the set \mathcal{D} has already been enumerated, and thus identified with a set $[0, M[$ of integers.) We let \mathcal{M} be the free \mathcal{A} -module generated by \mathcal{D} , where $\mathcal{A} = \mathbf{Z}[v, v^{-1}]$, and we write $q = v^2$. We replace the canonical basis $(T_\delta)_{\delta \in \mathcal{D}}$ of \mathcal{M} by $t_\delta = v^{-l(\delta)}T_\delta$, where $l : \mathcal{D} \rightarrow \mathbf{N}$ is the length function (also assumed to be tabulated), and similarly denote t_w the corresponding basis of the Hecke algebra \mathcal{H} of the complex Weyl group. We denote i the canonical involution on \mathcal{M} : the unique \mathcal{A} -antilinear involution such that

$$\begin{aligned} i(hm) &= i(h)i(m) && \text{for all } h \in \mathcal{H}, m \in \mathcal{M} \\ i(t_\delta) &= t_\delta + \sum_{\gamma < \delta} r(\gamma, \delta) t_\gamma && r(\gamma, \delta) \in \mathcal{A} \end{aligned}$$

The *Kazhdan–Lusztig basis* of \mathcal{M} is the unique basis $(c_\delta)_{\delta \in \mathcal{D}}$ such that :

$$\begin{aligned} i(c_\delta) &= c_\delta && \text{for all } \delta \in \mathcal{D} \\ c_\delta &= t_\delta + \sum_{\gamma < \delta} p(\gamma, \delta) t_\gamma && p(\gamma, \delta) \in v^{-1}\mathbf{Z}[v^{-1}] \end{aligned}$$

We denote $\mu(\gamma, \delta)$ the coefficient of v^{-1} in $p(\gamma, \delta)$.

Denote $\mathcal{M}_{\leq} = \bigoplus_{\delta \in \mathcal{D}} \mathbf{Z}[v^{-1}]t_{\delta}$, $\mathcal{M}_{<} = \bigoplus_{\delta \in \mathcal{D}} v^{-1}\mathbf{Z}[v^{-1}]t_{\delta}$. Then it is easy to show that if $m \in \mathcal{M}_{\leq}$ is such that $i(m) = m$, there are uniquely defined integers a_{δ} such that

$$m = \sum_{\delta \in \mathcal{D}} a_{\delta} c_{\delta}$$

Indeed, there are certainly uniquely defined *polynomials* $a_{\delta} \in \mathbf{Z}[v^{-1}]$ such that $m = \sum_{\delta \in \mathcal{D}} a_{\delta} c_{\delta}$, and then $i(m) = m$ is equivalent to $i(a_{\delta}) = a_{\delta}$ for all δ , which is equivalent to $a_{\delta} \in \mathbf{Z}$.

We denote $\tau(\delta)$ the descent set for δ (this may also be assumed to be tabulated.) In other words, $\tau(\delta)$ is the set of s that are either compact imaginary, complex such that $s \times \delta < \delta$, or real type I or type II for δ . The formulas for the action of the Hecke algebra are as follows :

(a) s is compact imaginary for δ :

$$\begin{aligned} T_s.T_{\delta} &= qT_{\delta} \\ t_s.t_{\delta} &= vt_{\delta} \\ c_s.t_{\delta} &= (v + v^{-1})t_{\delta} \end{aligned}$$

(b) s is noncompact imaginary type I for δ :

$$\begin{aligned} T_s.T_{\delta} &= T_{s \times \delta} + T_{\delta^s} \\ t_s.t_{\delta} &= v^{-1}t_{s \times \delta} + t_{\delta^s} \\ c_s.t_{\delta} &= v^{-1}t_{\delta} + v^{-1}t_{s \times \delta} + t_{\delta^s} \end{aligned}$$

(c) s is noncompact imaginary type II for δ :

$$\begin{aligned} T_s.T_{\delta} &= T_{\delta} + T_{\delta^s_+} + T_{\delta^s_-} \\ t_s.t_{\delta} &= v^{-1}t_{\delta} + t_{\delta^s_+} + t_{\delta^s_-} \\ c_s.t_{\delta} &= 2v^{-1}t_{\delta} + t_{\delta^s_+} + t_{\delta^s_-} \end{aligned}$$

(d) s is complex for δ , $s \in \tau(\delta)$:

$$\begin{aligned} T_s.T_{\delta} &= (q - 1)T_{\delta} + qT_{s \times \delta} \\ t_s.t_{\delta} &= (v - v^{-1})t_{\delta} + t_{s \times \delta} \\ c_s.t_{\delta} &= vt_{\delta} + t_{s \times \delta} \end{aligned}$$

(e) s is complex for δ , $s \notin \tau(\delta)$:

$$\begin{aligned} T_s.T_\delta &= T_{s \times \delta} \\ t_s.t_\delta &= t_{s \times \delta} \\ c_s.t_\delta &= v^{-1}t_\delta + t_{s \times \delta} \end{aligned}$$

(f) s is real type I for δ :

$$\begin{aligned} T_s.T_\delta &= (q-2)T_\delta + (q-1)(T_{\delta_s^+} + T_{\delta_s^-}) \\ t_s.t_\delta &= (v-2v^{-1})t_\delta + (1-v^{-2})(t_{\delta_s^+} + t_{\delta_s^-}) \\ c_s.t_\delta &= (v-v^{-1})t_\delta + (1-v^{-2})(t_{\delta_s^+} + t_{\delta_s^-}) \end{aligned}$$

(g) s is real type II for δ :

$$\begin{aligned} T_s.T_\delta &= (q-1)T_\delta - T_{s \times \delta} + (q-1)T_{\delta_s} \\ t_s.t_\delta &= (v-v^{-1})t_\delta - v^{-1}t_{s \times \delta} + (1-v^{-2})t_{\delta_s} \\ c_s.t_\delta &= vt_\delta - v^{-1}t_{s \times \delta} + (1-v^{-2})t_{\delta_s} \end{aligned}$$

(h) s is real for δ , $s \notin \tau(\delta)$:

$$\begin{aligned} T_s.T_\delta &= -T_\delta \\ t_s.t_\delta &= -v^{-1}t_\delta \\ c_s.t_\delta &= 0 \end{aligned}$$

2 Bruhat ordering

It turns out that there are in fact a number of orderings that one could use for the computation of Kazhdan–Lusztig polynomials. And in fact, what I am probably going to do is take the radical step of not using Bruhat order at all! (or rather, use the ordering where $\gamma \leq \delta$ iff $\gamma = \delta$ or $l(\gamma) < l(\delta)$.)

What we want from a “Bruhat ordering” is that (a) it is compatible with length (*i.e.*, weaker than the ordering described above, which we shall call the *length ordering*), and that (b) $p_{\gamma,\delta} \neq 0$ implies $\gamma \leq \delta$. Obviously, there is a strongest such ordering (the length ordering) and a weakest one (the one generated by the relations $p_{\gamma,\delta} \neq 0$.) Surprisingly, those two extremes do not seem to be all that far apart; even more so if instead of the weakest one, which seems rather hard to get at (is it even clear that it is graded?) we use one of the natural candidates, such as the one described in [2]. In

all the cases I computed, which included the large block for type E_7 , more than half the comparable pairs for the length ordering were also comparable for the Bruhat ordering in [2].

Recall the definition. Write $\delta \xrightarrow{s} \delta'$ if s is a strict descent for δ (*i.e.*, a complex descent, or real type I or II), and $\delta' = s \times \delta$ if s is a complex descent, $\delta' \in C_s(\delta)$ if s is real. Now let $\delta, \gamma \in \mathcal{D}$. Then we say that $\gamma < \delta$ iff there exists an $s \in S$ with $\delta \xrightarrow{s} \delta'$, and either $\gamma \leq \delta'$, or there is a $\gamma' < \delta'$ such that $\gamma \xrightarrow{s} \gamma'$. It should be fairly easy to show, pretty much as for the ordinary Bruhat order, that this ordering is graded, with degree being given by the length function. Then, in order to describe it, it suffices to give the covering relations.

As David pointed out to me, this ordering suffers from a lack of symmetry when passing over to the dual; one should want that the duality is order-reversing for the Bruhat ordering. As he explained, one solution would be to make the ordering even stronger, essentially by imposing that in the situation where $\delta \xrightarrow{s} \delta'$, $\gamma \xrightarrow{s} \gamma'$, one has $\gamma < \delta$ iff $\gamma' < \delta'$. It is not clear to me whether this definition is susceptible of a recursive construction; perhaps something of doing it as previously, and then applying symmetrizing passes until symmetry is achieved ?

Another solution, that would make the order weaker, would be to take the “product ordering” for the two projections on orbit sets, where on the orbit sets one takes the closure ordering as described combinatorially in Richardson and Springer [1]. I *think* this is an acceptable ordering, but this needs to be re-checked.

3 Kazhdan–Lusztig polynomials

Denote c_δ the Kazhdan-Lusztig basis. In order to compute the c_δ , we use two kinds of formulas :

I. Let $s \in \tau(\delta)$ (recall that this means that s is either imaginary compact, real in the domain of the inverse Cayley transform, or complex with $s \times \delta < \delta$.) then we have :

$$c_{s \cdot \delta} = (v + v^{-1}) c_\delta \tag{1}$$

This is Lemma 6.7 (b) in [2]. This yields the “easy case” induction formulas (second case type II in [2]) :

$$\begin{aligned}
p_{\gamma,\delta} &= v^{-1}p_{\gamma^s,\delta} \quad \text{for } s \text{ imaginary type I w.r.t. } \gamma \\
p_{\gamma,\delta} &= v^{-1}(p_{\gamma_+^s,\delta} + p_{\gamma_-^s,\delta}) \quad \text{for } s \text{ imaginary type II w.r.t. } \gamma \\
p_{\gamma,\delta} &= v^{-1}p_{s \times \gamma,\delta} \quad \text{for } s \text{ complex ascent w.r.t. } \gamma \\
p_{\gamma,\delta} &= 0 \quad \text{for } s \text{ real nonparity w.r.t. } \gamma
\end{aligned}$$

Indeed, we have :

$$\begin{aligned}
c_s \cdot c_\delta &= \sum_{\gamma \leq \delta} p_{\gamma,\delta} c_s \cdot t_\gamma = \sum_{\substack{\gamma < \delta \\ s \text{ compact}}} (v + v^{-1})p_{\gamma,\delta} t_\gamma \\
&+ \sum_{\substack{\gamma < \delta \\ s \text{ type I imaginary}}} [v^{-1}(p_{\gamma,\delta} + p_{s \times \gamma,\delta}) + (1 - v^{-2})p_{\gamma^s,\delta}] t_\gamma \\
&+ \sum_{\substack{\gamma < \delta \\ s \text{ type II imaginary}}} [2v^{-1}p_{\gamma,\delta} + (1 - v^{-2})(p_{\gamma_+^s,\delta} + p_{\gamma_-^s,\delta})] t_\gamma \\
&+ \sum_{\substack{\gamma < \delta \\ s \text{ complex descent}}} (vp_{\gamma,\delta} + p_{s \times \gamma,\delta}) t_\gamma + \sum_{\substack{\gamma < \delta \\ s \text{ complex ascent}}} (v^{-1}p_{\gamma,\delta} + p_{s \times \gamma,\delta}) t_\gamma \\
&+ \sum_{\substack{\gamma < \delta \\ s \text{ type I real}}} [(v - v^{-1})p_{\gamma,\delta} + p_{\gamma_+^s,\delta} + p_{\gamma_-^s,\delta}] t_\gamma \\
&+ \sum_{\substack{\gamma < \delta \\ s \text{ type II real}}} [vp_{\gamma,\delta} - v^{-1}p_{s \times \gamma,\delta} + p_{\gamma^s,\delta}] t_\gamma + \sum_{\substack{\gamma < \delta \\ s \text{ real nonparity}}} 0 \cdot t_\gamma
\end{aligned}$$

Consider now the four cases where s is an ascent for γ :

(a) s is noncompact type I imaginary with respect to γ . Then (1) yields :

$$v^{-1}(p_{\gamma,\delta} + p_{s \times \gamma,\delta}) + (1 - v^{-2})p_{\gamma^s,\delta} - (v + v^{-1})p_{\gamma,\delta} = 0 \quad (2)$$

But the same equation for γ^s yields :

$$(v - v^{-1})p_{\gamma^s,\delta} + p_{\gamma,\delta} + p_{s \times \gamma,\delta} - (v + v^{-1})p_{\gamma^s,\delta} = 0$$

Multiplying by v^{-1} , and using (2), we get

$$(v + v^{-1})(p_{\gamma,\delta} - v^{-1}p_{\gamma^s,\delta}) = 0$$

whence $p_{\gamma,\delta} = v^{-1}p_{\gamma^s,\delta}$.

(b) s is noncompact type II imaginary with respect to γ . Then (1) yields :

$$2v^{-1}p_{\gamma,\delta} + (1 - v^{-2})(p_{\gamma_+^s,\delta} + p_{\gamma_-^s,\delta}) - (v + v^{-1})p_{\gamma,\delta} = 0 \quad (3)$$

But the same equation for γ_+^s and γ_-^s yields :

$$\begin{aligned} vp_{\gamma_+^s,\delta} - v^{-1}p_{\gamma_-^s,\delta} + p_{\gamma,\delta} - (v + v^{-1})p_{\gamma_+^s,\delta} &= 0 \\ vp_{\gamma_-^s,\delta} - v^{-1}p_{\gamma_+^s,\delta} + p_{\gamma,\delta} - (v + v^{-1})p_{\gamma_-^s,\delta} &= 0 \end{aligned}$$

Multiplying by v^{-1} , adding up and combining with (3) yields :

$$(v + v^{-1})(p_{\gamma,\delta} - v^{-1}(p_{\gamma_+^s,\delta} + p_{\gamma_-^s,\delta})) = 0$$

whence $p_{\gamma,\delta} = v^{-1}(p_{\gamma_+^s,\delta} + p_{\gamma_-^s,\delta})$.

(c) s is complex for γ , $s \notin \tau(\gamma)$. Then (1) yields

$$v^{-1}p_{\gamma,\delta} + p_{s \times \gamma,\delta} - (v + v^{-1})p_{\gamma,\delta} = 0 \quad (4)$$

The same equation for $s \times \gamma$ yields :

$$vp_{s \times \gamma,\delta} + p_{\gamma,\delta} - (v + v^{-1})p_{s \times \gamma,\delta} = 0$$

Multiplying by v^{-1} and combining with (4) yields :

$$(v + v^{-1})(p_{\gamma,\delta} - v^{-1}p_{s \times \gamma,\delta}) = 0$$

whence $p_{\gamma,\delta} = v^{-1}p_{s \times \gamma,\delta}$.

(d) s is real for γ , $s \notin \tau(\gamma)$. Then (1) yields :

$$(v + v^{-1})p_{\gamma,\delta} = 0$$

whence $p_{\gamma,\delta} = 0$.

II. Let $s \notin \tau(\delta)$. Then there are four cases to consider :

(a) s is complex for δ , $s \times \delta > \delta$. Then :

$$c_s \cdot c_\delta = c_{s \times \delta} + \sum_{\substack{\zeta < \delta \\ s \in \tau(\zeta)}} \mu(\zeta, \delta) c_\zeta$$

(b) s is imaginary noncompact type I for δ . Then :

$$c_s \cdot c_\delta = c_{\delta s} + \sum_{\substack{\zeta < \delta \\ s \in \tau(\zeta)}} \mu(\zeta, \delta) c_\zeta$$

(c) s is imaginary noncompact type II for δ . Then :

$$c_s \cdot c_\delta = c_{\delta_+^s} + c_{\delta_-^s} + \sum_{\substack{\zeta < \delta \\ s \in \tau(\zeta)}} \mu(\zeta, \delta) c_\zeta$$

(d) s is real nonparity for δ . Then :

$$c_s \cdot c_\delta = \sum_{\substack{\zeta < \delta \\ s \in \tau(\zeta)}} \mu(\zeta, \delta) c_\zeta$$

III. Of course, in practice one would like to use the formulas in II for expressing c_δ in terms of lower parameters. This works as follows :

(a) s is complex for δ , $s \times \delta < \delta$. Then :

$$c_\delta = c_s \cdot c_{s \times \delta} - \sum_{\substack{\zeta < s \times \delta \\ s \in \tau(\zeta)}} \mu(\zeta, s \times \delta) c_\zeta$$

(b) s is real type I for δ . Then :

$$c_\delta = c_s \cdot c_{\delta_s^+} - \sum_{\substack{\zeta < \delta_s^+ \\ s \in \tau(\zeta)}} \mu(\zeta, \delta_s^+) c_\zeta$$

(here δ_s^- might have been used just as well instead of δ_s^+ .)

(c) s is real type II for δ . Then :

$$c_\delta + c_{s \times \delta} = c_s \cdot c_{\delta_s} - \sum_{\substack{\zeta < \delta_s \\ s \in \tau(\zeta)}} \mu(\zeta, \delta_s) c_\zeta$$

The last of these formulas is the one that doesn't directly give a recursion, and where the "structural fact" of Lemma 6.2 in [2] has to be used to show that the recursion can be gotten to go through. More specifically, if $G(\delta)$ is the connected component of δ in the graph where the edges are the $(\delta, s \times \delta)$ corresponding to generators in case (c) above, then for each fixed $\gamma < \delta$ there will be a $\delta' \in G(\delta)$ for which s is an ascent, allowing $p_{\gamma, \delta'}$ to be computed through one of the "easy" formulas obtained in I. Then the equations can be solved to get all the $p_{\gamma, \delta''}$, $\delta'' \in G(\delta)$.

In terms of polynomials, if we set δ' to be $s \times \delta$ in case (a), δ_s^+ in case (b), and δ_s in case (c), we get the following formulas for $\gamma < \delta$, where lhs denotes $p_{\gamma, \delta}$ in the first two cases, $p_{\gamma, \delta} + p_{\gamma, s \times \delta}$ in case (c)

(a) when s is a complex descent for γ :

$$\text{lhs} = v p_{\gamma, \delta'} + p_{s \times \gamma, \delta'} - \sum_{\gamma \leq \zeta < \delta'} \mu(\zeta, \delta') p_{\gamma, \zeta}$$

(b) s is real type I for γ . Then :

$$\text{lhs} = (v - v^{-1}) p_{\gamma, \delta'} + p_{\gamma_s^+, \delta'} + p_{\gamma_s^-, \delta'} - \sum_{\gamma \leq \zeta < \delta'} \mu(\zeta, \delta') p_{\gamma, \zeta}$$

(c) s is real type II for γ . Then :

$$\text{lhs} = v p_{\gamma, \delta'} - v^{-1} p_{s \times \gamma, \delta'} + p_{\gamma_s, \delta'} - \sum_{\gamma \leq \zeta < \delta'} \mu(\zeta, \delta') p_{\gamma, \zeta}$$

(d) s is imaginary noncompact for γ . Then :

$$\text{lhs} = (v + v^{-1}) p_{\gamma, \delta'} - \sum_{\gamma \leq \zeta < \delta'} \mu(\zeta, \delta') p_{\gamma, \zeta}$$

References

- [1] Roger W. Richardson and Tonny A. Springer. The Bruhat Order on Symmetric Varieties. *Geometriæ Dedicata*, **35**:389–436, 1990.
- [2] David A. Vogan, Jr. The Kazhdan-Lusztig conjecture for real reductive groups. In *Representation Theory of Reductive Groups (Park City, Utah, 1982)*, volume **40** of *Progress in Mathematics*, pages 223–264, Boston, Mass., 1983. Birkhäuser.